## Project 2 Signals

Deadline: October 10, Monday at 11.00 PM Minimum Requirement: 9 public tests

#### Test Breakdown

test breakdown: 12 public tests 1 release test 7 secret tests 5 points each except test\_built



# Project 2 requires you to implement signals and signal handlers.

#### Introduction to signals

- A signal is an inter-process communication mechanism that allows one process to invoke a signal-handler function in another process.
  - A process can send a signal to another process.
  - The process receiving the signal will, at some point, stop what it is doing, execute a signal-handler function, and then resume what it was doing.
- There are several signals that one process can send to another, each identified by a number.
  - Each process maintains a table of signal handlers (as function pointers), one for each signal the process can handle.
  - The signal number is used as an index into the table of signal handlers in the target process.
- When process A sends a signal to B, the kernel must create a new context that causes B to execute its signal handler, then return B to its original context.

#### Things to be implemented

<u>In syscall.c</u>

Sys\_Signal

Sys\_RegDeliver

Sys\_Kill

Sys\_ReturnSignal

Sys\_WaitNoPid

#### <u>In signal.c</u>

Check\_Pending\_Signal

Setup\_Frame

Complete\_Handler

## Sys\_Signal(struct Interrupt\_State \*state)

- A process calls "signal" to indicate what handler function should run when it receives a signal.
- This system call handler registers a signal handler for a given signal number.
- Recommend to check signal.h to know the structure
- Arguments: signal handler, signal number
  - state->ebx: pointer to handler function
  - state->ecx: signal number
- EINVALID ( can use IS\_SIGNUM to check if its a valid signal

#### Define & initialize a signal handlers table

- **SIG\_IGN**: tells the kernel the process wants to ignore the signal.
- **SIG\_DFL**: perform the default behavior of the process.
- Default behaviors:
  - SIGKILL, SIGPIPE, SIGUSER1, SIGUSER2: terminate
  - SIGCHLD: ignore

#### Include/geekos/signal.h

<pre>/* Signal number: #define SIGKILL</pre>		/* can't be	handled by use	ers */
<pre>#define SIGUSR1</pre>	2			
<pre>#define SIGUSR2</pre>	3			
<pre>#define SIGCHLD</pre>	4			
<pre>#define SIGALARM</pre>	5			
<pre>#define SIGPIPE</pre>	6			
/* The largest s: #define MAXSIG	ignal number su 6	pported */		

```
/* Macro to determine whether a number is a valid signal number */
#define IS SIGNUM(n) (((n) > 0) && ((n) <= MAXSIG))</pre>
```

```
/* Definition of a signal handler */
typedef void (*signal_handler) (int);
```

```
/* Default handlers */
#define SIG_DFL (signal_handler)1
#define SIG_IGN (signal_handler)2
```

#### **Return Signal and Trampoline**

- When a signal is sent to a process, the kernel arranges for the signal handler function to be called within the process.
- When the signal handler returns, control must pass back to the kernel, which arranges for the process to resume from wherever it was.
- To accomplish this, we define a system call **ReturnSignal** that is to be invoked by the process when its signal handler returns.
- However, we cannot count on every signal handler to call ReturnSignal explicitly.
  - Instead, the compiler will include a user-side function, called the "Trampoline," in every executable and set the return address of each signal handler to go to Trampoline.
- In userspace: calls Sys\_ReturnSignal syscall

#### Sys\_ReturnSignal(struct Interrupt\_State \*state)

- Complete signal handling for this process
- return state->eax;

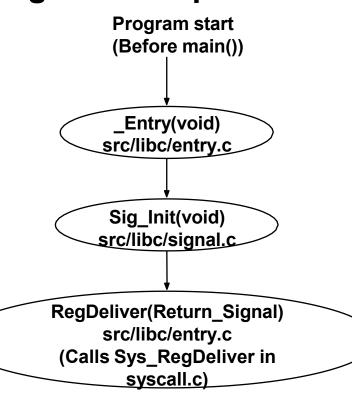
#### How does the kernel know where trampoline is?

• Sys\_Regdeliever is used to register the address at the beginning of each user program.

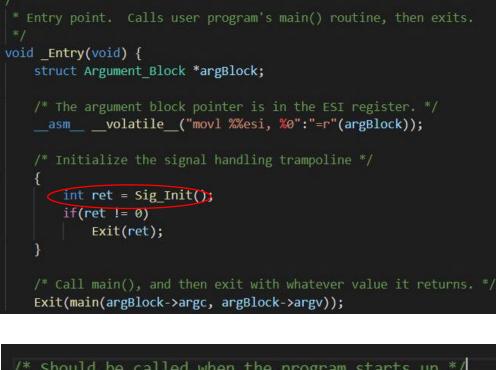
#### Sys\_RegDeliver(struct Interrupt\_State \*state)

- Must be invoked before a user process begins to execute.
- Provides the kernel with the address of the trampoline function, as described for Sys\_ReturnSignal.
- The user code that calls Sys\_RegDeliver and Sys\_ReturnSignal is already written for you in libc.
- Your task is :
  - To Store the values of the Return\_Signal(trampoline) address in User\_Context

#### Purpose: Register trampoline



Sys\_RegDeliver registers the Return\_Signal trampoline function for current process.



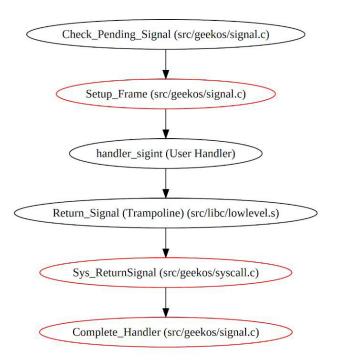
/\* Should be called when the program starts up \*/
int Sig\_Init(void) {
 return RegDeliver(Return\_Signal);

## Sys\_Kill(struct Interrupt\_State \*state)

- It will be used to send a signal to a certain process.
- If called with a different signal number, it should return EINVALID.
- Lookup\_Thread (src/geekos/kthread.c): get the target thread.
- Should be implemented as setting a flag in the process to which the signal is to be delivered, so that when the given process is about to start executing in user space again, rather than returning to where it left off, it will execute the appropriate signal handler instead.
- This is a kernel process; userspace isn't allowed to send signals to it. Set the appropriate error conditions.

#### Sys\_WaitNoPID(struct Interrupt\_State \*state)

- This call allows a parent to collect the exit status for a child that has terminated and become a zombie, without knowing its PID or going into a blocking wait.
- Argument: a pointer to an integer
- Reaps one zombie child process per call
- Places the exit status of the child process in the memory location the argument points to
- Returns the pid of the zombie.
- If no dead child process, return ENOZOMBIES



## Check\_Pending\_Signal(struct Kernel\_Thread \*kthread, struct Interrupt\_State \*state)

Check\_Pending\_Signal:

This routine is called by code in lowlevel.asm when a kernel thread is about to be context-switched in. It returns true if the following THREE conditions hold:

- 1. A signal is pending for that user process.
- The process is about to start executing in user space. This can be determined by checking the Interrupt\_State's CS register: if it is not equal to the kernel's CS register (see include/geekos/defs.h), then the process will return to user space.
- 3. The process is not currently handling another signal (recall that signal handling is non-reentrant).



Run every time when a process returns from kernel model to user mode

Basically checks if the process

- Has pending signal AND
- It's the right time to handle the signal
- Start executing in user space
- not handling another signal If

we passes these tests, then

- Iterate through all signal flags and set busy flag.
- Check for error conditions



- Called when switching to a user process, if Check\_Pending\_Signal returns true.
- Sets up a user processes user stack and kernel stack so that

(1) when the process returns to user mode, it will execute the correct signal handler, and

(2) when that handler completes, the process will return to the trampoline function so that it can go back to what it was doing.



- If the process has defined a signal handler for this signal, function Setup\_Frame will have to do the following:
  - 1. Choose the correct handler to invoke.
  - Acquire the pointer to the top of the *user stack*. This is available when you cast the saved Interrupt\_State (on the *kernel stack*) to a User\_Interrupt\_State.
  - 3. Push onto the *user stack* a snapshot of the interrupt state that is currently stored at the top of the *kernel stack*.
  - 4. Push onto the user stack the number of the signal being delivered.



- 5. Push onto the *user stack* the address of the trampoline (which invokes the Sys\_ReturnSignal system call handler). The trampoline was registered by the Sys\_RegDeliver system call handler, mentioned above.
- 6. Now that you have saved a copy of the *kernel stack*, change the original User\_Interrupt\_State such that
  - i. The *user stack* pointer is updated to reflect the changes made in steps 3--5.
  - ii. The saved program counter (eip) points to the signal handler.

#### Setup\_Frame(struct Kernel\_Thread \*kthread, struct Interrupt\_State \*state)

Find the registered handler pointer

•(cast to User\_Interrupt\_State to find user esp)

• Push user interrupt state to USER stack

• Push signal number to USER stack (argument to user defined signal handler pointer)

• Push trampoline to USER stack

•Jump to user defined signal handler pointer (state $\rightarrow$ eip)

	Increase
Previous User Stack Content	A
Interrupt State	8
Signal Number	•
Return Address (Trampoline)	S
Stack of User defined Signal Handler	
User Stack	

In Kernel Space

handle\_sigint()

User supplied function

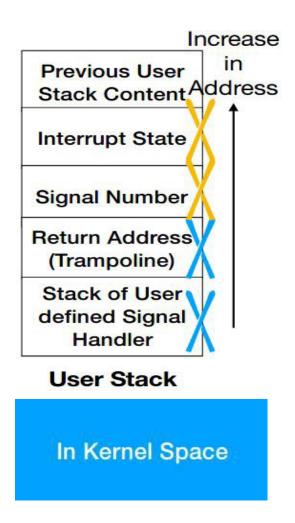
•Has function signature void signal\_handler(int signum);

Increase in **Previous User** Stack Content Address Interrupt State Signal Number **Return Address** (Trampoline) Stack of User defined Signal Handler User Stack

In User Space

#### Complete\_Handler(struct Kernel\_Thread \*kthread, struct Interrupt\_State \*state)

- This routine should be called when the Sys\_ReturnSignal call handler is invoked (after a signal handler has completed).
- It needs to restore back on the top of the kernel stack the snapshot of the interrupt state currently on the top of the user stack.
  - Reverse the operations in setup\_frame
  - Restore interrupt\_state on user stack to interrupt\_state on current thread
  - Remove signal number
- No need to pop the pushed return\_signal\_function here; it's been popped by the ret instruction at the end of the signal handler (to call the libc function Return\_Signal).



#### **Other Functions to Modify**

- 1. Sys\_Fork : copy signal handlers, trampoline
- 2. Pipe\_Write: send SIGPIPE
- 3. Exit: send SIGCHLD

- You have a lot of freedom for this project
- Feel free to add additional fields to existing structs
- •kernel\_thread for kernel space related data
- user\_context for user space related data

Demo: sigpipe.c

Remember the flow:

- Before main(): Sys\_RegDeliver -> Register trampoline function
- In main():
  - Check\_Pending\_Signal is always running
  - If user used Sys\_Signal to change a signal handler, when a signal is received (Check\_Pending\_Signal returns true):
    - i. Setup\_Frame (Kernel mode)
    - ii. User defined signal handler (User mode)
    - iii. Trampoline (User mode)
    - iv. Sys\_returnSignal (Kernel mode)
    - v. Complete\_Handler (Kernel mode)
    - vi. Back to user program (User mode)
  - If user didn't change the signal handler
    - i. Handle the signal in kernel (Terminate or ignore)
    - ii. Back to user program (User mode)

Demo: sigpipe.c

What this test is doing:

- The parent:
  - **Signal(sigpipe\_handler, SIGPIPE)**: It will register a signal handler sigpipe\_handler for SIGPIPE. So when the parent received SIGPIPE, it will execute the code in sigpipe\_handler.
  - Will try to read from the buffer. Then a read end will be closed.
  - Wait for the child to finish
- The child:
  - Close a read end.
  - Tries to write a lot of stuff to the buffer.
- Until this step, the two read ends should be closed. If any of the child or the parent tries to write to the buffer, a SIGPIPE will be raised before returning EPIPE.
- For the parent, SIGPIPE should invoke sigpipe\_handler. For the child, SIGPIPE should perform its default behavior.

